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| FINNEGAN, HENDERSON, FARABOW, GARRETT & DUNNER LLP 901 NEW YORK AVENUE, NW WASHINGTON, DC 20001-4413 | | | EXAMINER SANDERS, AARON J | |
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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

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|------------------------------|--------------------------------------|---|--|
| Office Action Summary | Application No. 10/526,750 | Applicant(s) PFERDEKAEMPER ET AL. | |
| | Examiner Aaron Sanders | Art Unit 2168 | |

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) ☒ Responsive to communication(s) filed on 27 July 2007.
- 2a) ☒ This action is **FINAL**. 2b) ☐ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) ☒ Claim(s) 1-9, 11, 12 and 15-31 is/are pending in the application.
- 4a) Of the above claim(s) _____ is/are withdrawn from consideration.
- 5) ☐ Claim(s) _____ is/are allowed.
- 6) ☒ Claim(s) 1-9, 11, 12 and 15-31 is/are rejected.
- 7) ☐ Claim(s) _____ is/are objected to.
- 8) ☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☒ The drawing(s) filed on 07 March 2007 is/are: a) ☐ accepted or b) ☒ objected to by the Examiner.
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some * c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
2. ☐ Certified copies of the priority documents have been received in Application No. _____.
3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).
- * See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

- | | |
|---|---|
| 1) <input type="checkbox"/> Notice of References Cited (PTO-892) | 4) <input type="checkbox"/> Interview Summary (PTO-413) Paper No(s)/Mail Date. _____ |
| 2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948) | 5) <input type="checkbox"/> Notice of Informal Patent Application |
| 3) <input checked="" type="checkbox"/> Information Disclosure Statement(s) (PTO/SB/08) Paper No(s)/Mail Date <u>06/19/2007</u> . | 6) <input type="checkbox"/> Other: _____ |

DETAILED ACTION

Response to Amendment

The amendment filed 27 July 2007 has been entered. Claims 1-9, 11-12, and 15-31 are pending. This action has been made FINAL as necessitated by amendment.

Drawings

The drawings are objected to under 37 CFR 1.83(a). The drawings must show every feature of the invention specified in the claims. Therefore, the method of claims 1 and 12 must be shown or the feature(s) canceled from the claim(s). No new matter should be entered.

Corrected drawing sheets in compliance with 37 CFR 1.121(d) are required in reply to the Office action to avoid abandonment of the application. Any amended replacement drawing sheet should include all of the figures appearing on the immediate prior version of the sheet, even if only one figure is being amended. The figure or figure number of an amended drawing should not be labeled as "amended." If a drawing figure is to be canceled, the appropriate figure must be removed from the replacement sheet, and where necessary, the remaining figures must be renumbered and appropriate changes made to the brief description of the several views of the drawings for consistency. Additional replacement sheets may be necessary to show the renumbering of the remaining figures. Each drawing sheet submitted after the filing date of an application must be labeled in the top margin as either "Replacement Sheet" or "New Sheet" pursuant to 37 CFR 1.121(d). If the changes are not accepted by the examiner, the applicant will be notified and informed of any required corrective action in the next Office action. The objection to the drawings will not be held in abeyance.

Claim Rejections - 35 USC § 112 First Paragraph

The following is a quotation of the first paragraph of 35 U.S.C. 112:

The specification shall contain a written description of the invention, and of the manner and process of making and using it, in such full, clear, concise, and exact terms as to enable any person skilled in the art to which it pertains, or with which it is most nearly connected, to make and use the same and shall set forth the best mode contemplated by the inventor of carrying out his invention.

Claims 1, 11, 12, and 23 are rejected under 35 U.S.C. 112, first paragraph, as failing to comply with the written description requirement. The claims contains subject matter which was not described in the specification in such a way as to reasonably convey to one skilled in the relevant art that the inventors, at the time the application was filed, had possession of the claimed invention. Specifically, the specification does not mention “displaying... whether the read and/or write access... can be performed” and “restricting a read and/or write access on the data object”.

Claim Rejections - 35 USC § 112 Second Paragraph

The following is a quotation of the second paragraph of 35 U.S.C. 112:

The specification shall conclude with one or more claims particularly pointing out and distinctly claiming the subject matter which the applicant regards as his invention.

Claims 1, 11, 12, and 23 are rejected under 35 U.S.C. 112, second paragraph, as being indefinite for failing to particularly point out and distinctly claim the subject matter which applicant regards as the invention.

Claims 1, 11, 12, and 23 are incomplete for omitting essential steps, such omission amounting to a gap between the steps. See MPEP § 2172.01. Specifically, the method does not perform any steps when the ID is not stored successfully.

Claim Rejections - 35 USC § 101

35 U.S.C. 101 reads as follows:

Whoever invents or discovers any new and useful process, machine, manufacture, or composition of matter, or any new and useful improvement thereof, may obtain a patent therefor, subject to the conditions and requirements of this title.

Claims 1-9, 11-12, and 15-31 are rejected under 35 U.S.C. 101 because the claimed invention is directed to non-statutory subject matter.

The disclosed subject matter lacks a practical application of a judicial exception (law of nature, abstract idea, naturally occurring article/phenomena) since it fails to produce a useful, concrete, and tangible result.

Specifically, the disclosed subject matter does not produce a tangible result because it fails to produce a result that is limited to having real world value rather than a result that may be interpreted to be abstract in nature as, for example, a thought, a computation, or manipulation of data. More specifically, the disclosed subject matter provides for performing read or write access on a data object, which may be a tangible result. However, since this does not occur in every situation and no other output is apparent in the disclosure, the invention lacks the necessary tangible result in every case. The produced result remains in the abstract and, thus, fails to achieve the required status of having real world value.

As per claims 11 and 23-31, the instant claims are directed to software *per se*. Independent claims 11 and 23 recite a computer program *per se* and functional descriptive material consisting of data structures and computer programs, which impart functionality when employed as a computer component. As such, the instant claims are not limited to statutory subject matter and are therefore non-statutory. See Diehr, 450 U.S. at 185-86, 209 USPQ at 8.

As per claims 12-22, according to the instant specification (see pg. 17, line 21), the computer readable medium includes propagation media, which are commonly carrier waves. As such, the instant claims are non-statutory.

Claim Rejections - 35 USC § 102

The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless –

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United States.

Claims 1-9, 11-12, and 15-31 are rejected under 35 U.S.C. 102(b) as being anticipated by U.S. 5,566,319 (a.k.a. Lenz).

As per claims 1-9, 11-12, and 15-31, Lenz teaches:

1. A method for accessing, in a computer system, a data object having an identifier (ID), comprising (*See e.g. col. 1, line 66 – col. 2, line 23, “a method of controlling access to data in storage which data and storage are shared by a plurality of processors”*):

storing the ID in a second lock object (*See e.g. Fig. 3, where the claimed “second lock object” is the referenced array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.*);

determining whether the ID was stored successfully, and upon a successful storage, determining, before accessing the data object, whether the ID is contained in a first lock object (*See e.g. Fig. 3 and col. 3, lines 1-10, “By using the search key ‘007’ during the execution of the lock instruction, the part of the lock file containing the control field for record 007 is addressed for writing” where the claimed “first lock object” is the referenced “Lock-File” 3-1*);

if the ID is contained in the first lock object, determining whether a link to a storage location is assigned to the ID in the first lock object (*See e.g. Fig. 3 where, see col. 1, line 66 – col. 2, line 23, “Each of the control fields is associated with a corresponding data address” and col. 3, lines 20-26, “As a result, by comparing the actual with the expected content of the status identification code associated with the control field, it is possible during writing to determine whether the write operation to be effected into the control field may be carried out”*), and

if the link is assigned to the ID, restricting a read and/or write access on the data object (*See e.g. Fig. 3 where, see col. 5, lines 51-55, “SKC could be 1: block not empty (at least one control field not empty), i.e. an access right may only be granted after detailed examination” where “SKC” means “status identification code”*), and

if the link is not assigned to the ID (*See e.g. Fig. 4 where, see col. 4, line 34 – col. 5, line 3, “access rights are withdrawn in steps 4-8, 4-9 and 4-10 (the unlock instruction is carried out), i.e. a new ZVI is entered in the lock file” where “ZVI” means “access administration information”*), deleting the ID from the first lock object (*See e.g. Fig. 4 where, see col. 4, line 34 – col. 5, line 3, “the status identification code is updated accordingly in step 4-10, for example, to the effect that the control field is now empty and may be overwritten without having to be read later on”*) and performing the read and/or write access on the data object (*See e.g. Fig. 3 where, see col. 5, lines 49-50, “SKC could be 0: block empty (all control fields empty), i.e. an access right may be granted immediately”*), and

if the ID is not contained in the first lock object, performing the read and/or write access on the data object (*See e.g. Fig. 3 and col. 3, lines 51-62, "If the examination of the status identification code SKC shows that the write operation into the control field for account record 007, refer to 3-2 for instance, may be executed, this write operation is carried out and the status identification code SKC is updated according to the write request"*); and

displaying to a user, on an output display, whether the read and/or write access on the data object is restricted or can be performed (*See e.g. Fig. 6 where, see col. 5, line 60 – col. 6, line 7, "FIG. 6 is a schematic time diagram illustrating how two processors at different times access a record addressable by the key '007' for executing a read instruction"*).

2. The method of claim 1, wherein the first lock object is a file stored on a nonvolatile storage means (*See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, "Each of the processors has a local main storage. The shared storage is located outside of the main storage and stores a lock file"*).

3. The method of claim 1, wherein the first lock object comprises a table having a first column for the ID and a second column for the link of the ID to a storage location (*See e.g. Fig. 3, "Lock-File" 3-1*).

4. The method of claim 1, wherein each data object comprises one or more fields of one or more tables and wherein the ID comprises one or more key fields of the one or more tables (*See e.g. Fig. 5*).

5. The method of claim 4, wherein the link is a filename or a link to a file (*See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “Each of the control fields is associated with a corresponding data address”*).

6. The method of claim 1, wherein the first lock object is created by a data moving process (*See e.g. col. 1, line 66 – col. 2, line 23, “a method of controlling access to data in storage which data and storage are shared by a plurality of processors”*).

7. The method of claim 1, wherein the second lock object is stored in a volatile storage means (*See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “Each of the processors has a local main storage. The shared storage is located outside of the main storage and stores a lock file”*).

8. The method of claim 1, wherein the second lock object is a data array (*See e.g. Fig. 3, the array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.*).

9. The method of claim 8 wherein the data array is one dimensional (*See e.g. Fig. 3, the array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.*).

10. (Canceled)

11. A computer system for processing data, comprising:

memory means for storing program instructions (*See e.g. Fig. 3, “Local Main Storage 1” 3-6*);

input means for entering data (*See e.g. Fig. 3, “Writing”*);

storage means for storing data (*See e.g. Fig. 3, shared storage 3-9*);

a processor responsive to the program instructions, wherein the program instructions comprise program code means for performing a method for accessing a data object having an identifier (ID), the method comprising (*See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “a*

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method of controlling access to data in storage which data and storage are shared by a plurality of processors”):

storing the ID in a second lock object (See e.g. Fig. 3, where the claimed “second lock object” is the referenced array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.);

determining whether the ID was stored successfully, and upon a successful storage, determining, before accessing the data object, whether the ID is contained in a first lock object (See e.g. Fig. 3 and col. 3, lines 1-10, “By using the search key ‘007’ during the execution of the lock instruction, the part of the lock file containing the control field for record 007 is addressed for writing” where the claimed “first lock object” is the referenced “Lock-File” 3-1);

if the ID is contained in the first lock object, determining whether a link to a storage location is assigned to the ID in the first lock object (See e.g. Fig. 3 where, see col. 1, line 66 – col. 2, line 23, “Each of the control fields is associated with a corresponding data address” and col. 3, lines 20-26, “As a result, by comparing the actual with the expected content of the status identification code associated with the control field, it is possible during writing to determine whether the write operation to be effected into the control field may be carried out”), and

if the link is assigned to the ID, restricting a read and/or write access on the data object (See e.g. Fig. 3 where, see col. 5, lines 51-55, “SKC could be 1: block not empty (at least one control field not empty), i.e. an access right may only be granted after detailed examination”), and

if the link is not assigned to the ID (See e.g. Fig. 4 where, see col. 4, line 34 – col. 5, line 3, “access rights are withdrawn in steps 4-8, 4-9 and 4-10 (the

unlock instruction is carried out), i.e. a new ZVI is entered in the lock file”), deleting the ID from the first lock object (See e.g. Fig. 4 where, see col. 4, line 34 – col. 5, line 3, “the status identification code is updated accordingly in step 4-10, for example, to the effect that the control field is now empty and may be overwritten without having to be read later on”) and performing the read and/or write access on the data object (See e.g. Fig. 3 where, see col. 5, lines 49-50, “SKC could be 0: block empty (all control fields empty), i.e. an access right may be granted immediately”), and

if the ID is not contained in the first lock object, performing the read and/or write access on the data object (See e.g. Fig. 3 and col. 3, lines 51-62, “If the examination of the status identification code SKC shows that the write operation into the control field for account record 007, refer to 3-2 for instance, may be executed, this write operation is carried out and the status identification code SKC is updated according to the write request”); and

displaying to a user, on an output display, whether the read and/or write access on the data object is restricted or can be performed (See e.g. Fig. 6 where, see col. 5, line 60 – col. 6, line 7, “FIG. 6 is a schematic time diagram illustrating how two processors at different times access a record addressable by the key ‘007’ for executing a read instruction”).

12. A computer readable medium comprising instructions for performing a method for accessing a data object having an identifier in a computer system, the method comprising (See e.g. col. 1, line 66 – col. 2, line 23, “a method of controlling access to data in storage which data and storage are shared by a plurality of processors”):

storing the ID in a second lock object (See e.g. Fig. 3, where the claimed “second lock object” is the referenced array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.); and

determining whether the ID was stored successfully, and upon a successful storage, determining, before accessing the data object, whether the ID is contained in a first lock object (See e.g. Fig. 3 and col. 3, lines 1-10, “By using the search key ‘007’ during the execution of the lock instruction, the part of the lock file containing the control field for record 007 is addressed for writing” where the claimed “first lock object” is the referenced “Lock-File” 3-1);

if the ID is contained in the first lock object, determining whether a link to a storage location is assigned to the ID in the first lock object (See e.g. Fig. 3 where, see col. 1, line 66 – col. 2, line 23, “Each of the control fields is associated with a corresponding data address” and col. 3, lines 20-26, “As a result, by comparing the actual with the expected content of the status identification code associated with the control field, it is possible during writing to determine whether the write operation to be effected into the control field may be carried out”), and

if the link is assigned to the ID, restricting a read and/or write access on the data object (See e.g. Fig. 3 where, see col. 5, lines 51-55, “SKC could be 1: block not empty (at least one control field not empty), i.e. an access right may only be granted after detailed examination”), and

if the link is not assigned to the ID (See e.g. Fig. 4 where, see col. 4, line 34 – col. 5, line 3, “access rights are withdrawn in steps 4-8, 4-9 and 4-10 (the unlock instruction is carried out), i.e. a new ZVI is entered in the lock file”), deleting the ID from the first lock object (See e.g. Fig. 4 where, see col. 4, line 34

– col. 5, line 3, “the status identification code is updated accordingly in step 4-10, for example, to the effect that the control field is now empty and may be overwritten without having to be read later on”) and performing the read and/or write access on the data object (See e.g. Fig. 3 where, see col. 5, lines 49-50, “SKC could be 0: block empty (all control fields empty), i.e. an access right may be granted immediately”), and

if the ID is not contained in the first lock object, performing the read and/or write access on the data object (See e.g. Fig. 3 and col. 3, lines 51-62, “If the examination of the status identification code SKC shows that the write operation into the control field for account record 007, refer to 3-2 for instance, may be executed, this write operation is carried out and the status identification code SKC is updated according to the write request”); and

displaying to a user, on an output display, whether the read and/or write access on the data object is restricted or can be performed (See e.g. Fig. 6 where, see col. 5, line 60 – col. 6, line 7, “FIG. 6 is a schematic time diagram illustrating how two processors at different times access a record addressable by the key ‘007’ for executing a read instruction”).

13. (Canceled)

14. (Canceled)

15. The computer readable medium of claim 12, wherein the first lock object is a file stored on a nonvolatile storage means (See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “Each of the processors has a local main storage. The shared storage is located outside of the main storage and stores a lock file”).

16. The computer readable medium of claim 12, wherein the first lock object comprises a table having a first column for the ID and a second column for the link of the ID to a storage location (*See e.g. Fig. 3, "Lock-File" 3-1*).

17. The computer readable medium of claim 12, wherein each data object comprises one or more fields of one or more tables and wherein the ID comprises one or more key fields of the one or more tables (*See e.g. Fig. 5*).

18. The computer readable medium of claim 12, wherein the link is a filename or a link to a file (*See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, "Each of the control fields is associated with a corresponding data address"*).

19. The computer readable medium of claim 12, wherein the first lock object is created by a data moving process (*See e.g. col. 1, line 66 – col. 2, line 23, "a method of controlling access to data in storage which data and storage are shared by a plurality of processors"*).

20. The computer readable medium of claim 12, wherein the second lock object is stored in a volatile storage means (*See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, "Each of the processors has a local main storage. The shared storage is located outside of the main storage and stores a lock file"*).

21. The computer readable medium of claim 12, wherein the second lock object is a data array (*See e.g. Fig. 3, the array of "Shared Data Record[s]" for "006", "007", "008", etc.*).

22. The computer readable medium of claim 21, wherein the data array is one dimensional (*See e.g. Fig. 3, the array of "Shared Data Record[s]" for "006", "007", "008", etc.*).

23. A computer system for processing data, comprising (*See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “a method of controlling access to data in storage which data and storage are shared by a plurality of processors”*):

means for storing an identifier (ID) in a second lock object (*See e.g. Fig. 3, where the claimed “second lock object” is the referenced array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.*); and

means for determining whether the ID was stored successfully, and upon a successful storage, determining, before accessing the data object, whether the ID is contained in a first lock object (*See e.g. Fig. 3 and col. 3, lines 1-10, “By using the search key ‘007’ during the execution of the lock instruction, the part of the lock file containing the control field for record 007 is addressed for writing” where the claimed “first lock object” is the referenced “Lock-File” 3-1*);

if the ID is contained in the first lock object, determining whether a link to a storage location is assigned to the ID in the first lock object (*See e.g. Fig. 3 where, see col. 1, line 66 – col. 2, line 23, “Each of the control fields is associated with a corresponding data address” and col. 3, lines 20-26, “As a result, by comparing the actual with the expected content of the status identification code associated with the control field, it is possible during writing to determine whether the write operation to be effected into the control field may be carried out”*), and

if the link is assigned to the ID, restricting a read and/or write access on the data object (*See e.g. Fig. 3 where, see col. 5, lines 51-55, “SKC could be 1: block not empty (at least one control field not empty), i.e. an access right may only be granted after detailed examination”*), and

if the link is not assigned to the ID (See e.g. Fig. 4 where, see col. 4, line 34 – col. 5, line 3, “access rights are withdrawn in steps 4-8, 4-9 and 4-10 (the unlock instruction is carried out), i.e. a new ZVI is entered in the lock file”), deleting the ID from the first lock object (See e.g. Fig. 4 where, see col. 4, line 34 – col. 5, line 3, “the status identification code is updated accordingly in step 4-10, for example, to the effect that the control field is now empty and may be overwritten without having to be read later on”) and performing the read and/or write access on the data object (See e.g. Fig. 3 where, see col. 5, lines 49-50, “SKC could be 0: block empty (all control fields empty), i.e. an access right may be granted immediately”), and

if the ID is not contained in the first lock object, performing the read and/or write access on the data object (See e.g. Fig. 3 and col. 3, lines 51-62, “If the examination of the status identification code SKC shows that the write operation into the control field for account record 007, refer to 3-2 for instance, may be executed, this write operation is carried out and the status identification code SKC is updated according to the write request”); and

means for displaying to a user, on an output display, whether the read and/or write access on the data object is restricted or can be performed (See e.g. Fig. 6 where, see col. 5, line 60 – col. 6, line 7, “FIG. 6 is a schematic time diagram illustrating how two processors at different times access a record addressable by the key ‘007’ for executing a read instruction”).

24. The computer system of claim 23, wherein first lock object is a file stored on a nonvolatile storage means (See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “Each of the

processors has a local main storage. The shared storage is located outside of the main storage and stores a lock file”).

25. The computer system of claim 23, wherein the first lock object comprises a table having a first column for the ID and a second column for the link of the ID to a storage location (See e.g. Fig. 3, “Lock-File” 3-1).

26. The computer system of claim 23, wherein each data object comprises one or more fields of one or more tables and wherein the ID comprises one or more key fields of the one or more tables (See e.g. Fig. 5).

27. The computer system of claim 23, wherein the link is a filename or a link to a file (See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “Each of the control fields is associated with a corresponding data address”).

28. The computer system of claim 23, wherein the first lock object is created by a data moving process (See e.g. col. 1, line 66 – col. 2, line 23, “a method of controlling access to data in storage which data and storage are shared by a plurality of processors”).

29. The computer system of claim 23, wherein the second lock object is stored in a volatile storage means (See e.g. Fig. 3 and col. 1, line 66 – col. 2, line 23, “Each of the processors has a local main storage. The shared storage is located outside of the main storage and stores a lock file”).

30. The computer system of claim 23, wherein the second lock object is a data array (See e.g. Fig. 3, the array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.).

31. The computer system of claim 30, wherein the data array is one dimensional (See e.g. Fig. 3, the array of “Shared Data Record[s]” for “006”, “007”, “008”, etc.).

Response to Arguments

As per Applicant's argument that claims 1, 11, 12, and 23 are not indefinite under 35 U.S.C. 112 second paragraph, the Examiner respectfully disagrees. Specifically, the method does not perform any steps when the ID is not stored successfully. The specification may disclose the missing steps, but those steps are not read into the claim limitations. Thus, the claims are indefinite because the steps performed when the ID is not stored successfully are unknown.

As per Applicant's argument that claims 1-9, 11-12, and 15-31 are statutory under 35 U.S.C. 101, the Examiner respectfully disagrees. As per independent claims 1, 11, 12, and 23, the step of displaying may be a tangible result, but since it does not appear in the specification it cannot be considered. Further, performing read or write access on a data object may be a tangible result if it is the final result of the method as a whole and occurs in every case. This does not appear to be the case, and no other output is apparent in the disclosure.

As per system claims 11 and 23-31, the claimed "means" are not necessarily hardware. Generally a system requires one or more processors coupled to a storage device containing instructions that when executed perform the method of...

As per claims 12-22, according to the instant specification (see pg. 17, line 21), the computer readable medium includes propagation media, which are commonly carrier waves.

As per Applicant's argument that Lenz does not disclose, "if the link is not assigned to the ID, deleting the ID from the first lock object" in claim 1, the Examiner respectfully disagrees. The Examiner cited Fig. 4 where, see col. 4, line 34 – col. 5, line 3, "the status identification

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code is updated accordingly in step 4-10, for example, to the effect that the control field is now empty and may be overwritten without having to be read later on” and col. 4, line 34 – col. 5, line 3, “access rights are withdrawn in steps 4-8, 4-9 and 4-10 (the unlock instruction is carried out), i.e. a new ZVI is entered in the lock file” as disclosing the claimed feature.

Conclusion

Applicant’s amendment necessitated the new ground(s) of rejection presented in this Office action. Accordingly, **THIS ACTION IS MADE FINAL**. See MPEP § 706.07(a). Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the date of this final action.

Any inquiry concerning this communication or earlier communications from the Examiner should be directed to Aaron Sanders whose telephone number is 571-270-1016. The Examiner can normally be reached on M-Th 8:00a-5:00p.

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If attempts to reach the Examiner by telephone are unsuccessful, the Examiner's supervisor, Tim Vo can be reached on 571-272-3642. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

/AJS/
Aaron J. Sanders
Examiner
18 September 2007

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